# Building Melodic Feature Knowledge of Gamelan Music Using Apriori Based on Functions in Sequence (AFiS) Algorithm 

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#### Abstract

Gamelan is a traditional music ensemble from Java, Indonesia, whose melody has characteristics that make the melodic sound of gamelan music easy to recognize. This research aims at building melodic feature knowledge of gamelan music in terms of note sequences rules. The algorithm called AFiS (Apriori based on Functions in Sequence) was also introduced to produce rules by mining the frequent value of note sequences. The basic idea of the AFiS algorithm is to define functions in a sequence, and then to chain the functions based on its position order to identify the support value for each function. The implementation of AFiS algorithm is aimed to define rules of gamelan music melodic feature in terms of ideal note sequences for composition. The evaluation of the accuracy of the note sequences rules is conducted by developing a recommendation system using rules defined in this research. The program is expected to answer correctly to some notes randomly deleted from the sequences. The result shows that the accuracy of the knowledge, and that the note sequences rules of gamelan music based on the correct answer is up to $86.5 \%$. Another evaluation is to find whether the different answers given by the program are accepted as alternative notes to the original notes. This evaluation involved 4 human experts to describe their acceptance of the alternative notes based on the different answers. The result shows that the different notes in 4 of 5 gendings are accepted by the experts as alternative notes. Copyright © 2016 Praise Worthy Prize S.r.l. - All rights reserved.


Keywords: AFiS Algorithm, Sequential Pattern Mining, Gamelan Music

|  | Nomenclature |
| :--- | :--- |
| A-B-C-D | A concept of gatra |
| D | Database |
| F | Series of functions |
| I | Collection of items |
| i | items |
| LN | Frequent sequence |
| Minsup | Minimum support |
| P | Data partition |
| PID | Partition ID |
| PI | Itemsets in a partition |
| S | Sequence |
| SID | Sequence ID |
| TC | Total number of candidate |
| TF | Total number of function |
| TS | Total number of sequence |
| TSI | Total number of itemsets in a sequence |

## I. Introduction

The use of artificial intelligence in music composition process is known as algorithmic composition, where a certain algorithm is used to automatically create a music composition [1] [2] [3]. Algorithmic composition has been developed since mid-year 1950, when Lejaren A. Hiller implemented a computer-generated composition by designing a computer to generate number sequences
randomly representing notes, rhythm, tempo, and others, to create Illiac Suite [4]. Xenakis was another pioneer work which used stochastic algorithm and Markov model to generate basic materials for music composition [5]. An expert systems approach has been used in CHORAL to harmonize four-part chorales in the style of J.S. Bach by defining 270 rules for chord skeleton, individual melodic lines, and others [6]. Another example of the use of rulebased for the music composition is represented by Baroc Harmony which added a pre-defined melody to control the search space [7]. The grammar approach was used by [8] to develop Improve Generator, a system that learns the pattern of percussion in live-streaming, and generates an accompaniment track which is improvised in realtime.

Considering the growth of the use of the algorithmic composition approach in many musical genres, the present paper uses the approach for a traditional music orchestra called karawitan. Karawitan known also as gamelan music is a music genre originally coming from Java, Indonesia. Karawitan uses gamelan as orchestra music instruments, and gending as song. The final goal of this research is to develop a rule-based expert system for the automatic gamelan music composition. One of the main parts of the system developed in this research is to build a knowledge base for the system to arrange a note sequence that could fit the melodic feature of gamelan music. Therefore, this paper proposes a model to build
the melodic feature knowledge of gamelan music using the sequential pattern mining technique implemented by identifying the association among notes in a sequence. It also introduces an algorithm, called AFiS (Apriori based on Functions in Sequence), as a sequence data mining algorithm used to analyze the sequential patterns of the gamelan song notes.

## II. Related Researches

One of the researches in algorithmic composition is a phenomenal work conducted by [5], which used an expert systems approach to develop a system called CHORAL. The system aimed at harmonizing four-part chorales in the style of Johann Sebastian Bach. The knowledge models of the CHORAL system used 270 rules consisting of: the chord skeleton which controls the rules to generate a new key; the fill-in which observes the chorale as four interacting automata, and read the output of the chord skeleton for enumerating the possible inessential note patterns; the time-slice which controls the constraint about the consecutive octaves and fifths; the melodic string which observes the sequence of individual notes of the different voices from a purely melodic point of view; the merged melodic string which merges the repeated adjacent pitches into a single note; The Schenkerian analysis based on formal theory of hierarchical voice leading to controls the rules of the possible parser actions, a constraint about the key of the chorale, and a heuristic for proper recognition of a Schenkerian $D-C-B-C$ ending pattern.

An application based on L-system for music generation was developed by [9] by rendering the contour of music using F which denotes the generation of a note, f, a pause, and + versus -, the increase, versus the decrease of the pitch of the played note. In the conclusion of the experiment, [9] stated that L-system could not be a first choice for the direct composition of traditional music, but could be ideal for the composition of interactive music.

A class of generative grammars called Probabilistic Temporal Graph Grammars (PTGG) implemented in Haskell was proposed by [10]. This class aimed at handling the temporal aspect of the music in a way that could retain a coherent metrical structure. The concept of PTGG is described as: The grammar generates some sequences of duration-parameterized abstract chord. Chords function as both terminals and non-terminals; ignoring the duration in which the context of a chord does not effect the productions that may be applied to it. Rules are functions on the duration of their input symbol. A system called Kulitta proposed by [11] used a Haskellbased framework for automated and algorithmic music composition. The modules of Kulitta are divided into 3 main categories: abstract/structural generation which contains a collection of schenkerian-inspired models and algorithms for iteratively generating harmonic structure or other abstract; musical interpretation which contains mathematical models and constraint satisfaction
algorithms for turning the abstract musical features into more detailed music; learning which contains offline learning algorithms to derive the production probabilities for musical grammars. Kulitta features a category of grammars called Probabilistic Temporal Graph Grammars (PTGGs) which parameterize the harmonic symbols in duration, and the rules are the functions on that duration.

Other examples of works in automatic music generation are focused on more specific tasks with reductionist approach, such as in [12] [13]. The reorchestration of Beethoven's Ode to Joy, the European anthem, in seven styles of Bach chorales, bossa nova, jazz, lounge, penny lane, chi mai, prayer in C, was served using machine learning approach [12]. There were some different approaches in generating different style of music, for instance, in Bach chorales style, the sequence to be generated by a probability distribution is represented using the max entropy model, while the fioriture-based Markov model is used in jazz style. A Natural Language Processing technique is used by [13] for the automatic reduction of melodies using the Probabilistic Context-Free Grammar (PCFG). The PCFG for melodic reduction tasks is trained using supervised learning, and the three analyses provided by the Generative Theory of Tonal Music (GTTM) dataset is used for the learning process. The performance of the melodic reduction includes a melodic identification and similarity, an efficient storage of melodies, the automatic composition, the variation matching, and the automatic harmonic analysis. The natural language text was used to generate music automatically by [14]. A sequence of notes, as well as the rhythm, was generated using text input mapped using vowels in conjunction with the word's part of speech (POS).

The learning technique, probability and statistical analysis were used to generate new music, never heard before, by [15]. Note sequences and other musical parameters such as note length, pitch, accidentals, modification (intensity, speed), and note sequence repetition density were used for the preparation of a probability table that could generate new music. There are three phases to generate new music: Music database creation that stores note identifier of the current note, the note identifier of the previous note, current sequence identifier; Learning and analysis phase based on thematic piece (i.e. rock, classical, instrumental, etc) using statistical analysis to produce a generalized rule-based matrix of note and sequence identifiers for creating new pieces; Music generation phase that generate new music based on the matrix created during statistical analysis.

In gamelan music, there are skeleton notes called balungan. The gamelan composer usually creates the skeleton notes firstly, before completing the composition. The note sequence in a gending is arranged based on gatra. Gatra is the smallest unit that consists of 4 sabetan balungan (beats). The research in gamelan music conducted by [16] used a grammar approach to define the rules of gending (gamelan song) composition in a style
of gamelan music called srepegan, by identifying the contour of note sequences. The pattern of srepegan is formulated based on $\operatorname{gatra}(s)$, which are the smallest unit consisting of 4 beats, and is served in term of a contour which represents the notes scale, where the frequent patterns are identified by higher or lower than the previous note. Another approach was used by [17] by building a frame work of quasi-linguistic approach to describe the structure and content of a style of gamelan music called gending lampah by defining rules of composition according to gending lampah. The grammar consists of the base rules, which are common rules that control the structure of gending lampah, the contour assignment rules which control the pitch assignment, the restriction rules which restrict the range of choice of the transformation rules, the transformation rules which control the choice of the notes, and their arrangement, the derivation rules which set up to generate a gending lampah, the gong assignment rules which controls the assignment of gong which is a gamelan instruments. Both these researches used frequent pattern mining to define the rules of a particular style of gamelan music. Although the subject of research was different in style or genre, the contour pattern proposed by [16], and quasilinguistic approach proposed by [17] involved the melodic feature of gamelan music as a part of the component controlled by the rules. Instead of using frequent pattern mining, this research preferred to use a sequential pattern mining in building melodic feature knowledge of gamelan music. This technique was expected to result in rules which could more generally be applied to gamelan music.

Trends of sequential patterns mining for big data have grown rapidly, such as in the works by [18] which developed Peekquence; a sequential pattern mining approach to explore event sequences in big data, which used users relevant metrics to sort patterns, to integrate patterns with patient time lines, and to overview and summarize the most common events in patterns. The traditional approaches of sequential patterns mining are considered as a design which does not support a massive amount of data. However traditional approaches of sequential pattern mining still need continuous research.

Kamber et al in [19] stated that sequential patterns are multidimensional association rules; therefore there are different specialized ways to find sequential patterns. Further, it is mentioned that some different approaches in mining sequential patterns are Apriori-based algorithm GSP, SPAM, projection-based FreeSpan, PSPM, vertical data format based algorithm SPADE, and pattern growth based approach in UDDAG.

The sequential patterns mining algorithm is categorized based on its structure by the type of Aprioribased and Pattern Growth-based [20] [21] [22] dataset. Apriori-based algorithm, such as AprioriAll, AprioriSome, DynamicSome, GSP (Generalized Sequential Patterns), SPADE (Sequential Pattern Discovery using Equivalence classes), SPAM (Sequential PAttern Mining), has typically been used to
discover intra-transaction associations and generate rules about the discovered associations, the example algorithm; Pattern Growth Algorithms, such as FreeSpan, and PrefixSpan, remove the process of candidate generation and prune steps that occur in the Apriori-type algorithms. The pattern growth algorithms can be faster when given large volumes of data; Temporal sequences, such as WINEPI, MINEPI, PROWL, allows the data used for sequence mining not to be limited to data stored in overtly temporal or longitudinally maintained datasets [21].

P-Prefix Span (Percussive-Prefix Sequential pattern) algorithm is used by [23] to integrate the genetic information and discover complex diseases such as Type2 Diabetes Mellitus. The P-Prefix Span algorithm works by discovering the length of sequential pattern and counting the support for all gene sequences to generate interesting patterns which satisfy the conditions. The DBP-SPAM algorithm which employs direct bit position manipulation technique is proposed by [24]. The frequent itemsets are discovered by bit position pruning technique.

There are three features in the taxonomy of Aprioribased algorithm defined by [25]: Breadth-first search, where the algorithm construct all $k$-sequences together in each $k_{\mathrm{th}}$ iteration of the algorithm as they traverse the search space; Generate-and test, in which the pattern is simply grown one item at a time and tested against the minimum support; Multiple scans of the database that scan the original database to ascertain whether a long list of generated candidate sequences is frequent or not.

In this research a new algorithm for sequential pattern mining is proposed, in which the Apriori-based is used to construct procedures of the algorithm.

## III. Proposed Work

A sequential pattern mining technique is used to build melodic feature knowledge of gamelan music which is able to represent the $A-B-C-D$ concept of gamelan music. We use the technique to formulate an appropriate note sequence by finding the frequent notes of a function that implies the existence of the frequent notes of the following function. The use of functions in a sequence is considered for mining a sequential pattern. In this research, the algorithm called Apriori based on Function in Sequence (AFiS) is proposed.

The AFiS algorithm works based on functions in a sequence. A function contains an item based on its order. The functions are then chained in terms of sequential patterns. Finally, the rules are defined based on frequent sequences. The AFiS algorithm has seven phases in mining sequence data: function definition, data partition, sequential pattern creation, candidate selection, support counting, prune phase, and production rules.

In this explanation, there is database $D$, which contains sequences ( S ), where each $S$ consists of an ordered list of itemset < $\mathrm{s} 1, \mathrm{~s} 2, \ldots, \mathrm{sn}>$, and the $s$ sequence contains items. We denote the collection of items by $I$,
where $I=\left\{i_{1}, i_{2}, . ., i_{n}\right\}$. We use a sequence database below, as an example in explaining the implementation of AFiS algorithm (Fig. 1).

| SID | Transaction |
| :---: | :--- |
| 01 | $\mathrm{c}, \mathrm{f}, \mathrm{a}, \mathrm{b}, \mathrm{e}, \mathrm{c}, \mathrm{b}, \mathrm{a}, \mathrm{c}, \mathrm{d}, \mathrm{a}$ |
| 02 | $\mathrm{f}, \mathrm{b}, \mathrm{d}, \mathrm{a}, \mathrm{c}, \mathrm{a}, \mathrm{c}, \mathrm{f}$ |
| 03 | $\mathrm{~b}, \mathrm{e}, \mathrm{d}, \mathrm{e}, \mathrm{f}, \mathrm{a}$ |

Fig. 1. Sequence database
Based on the explanation above, then:

$$
\begin{aligned}
& \mathrm{I}=\{\mathrm{a}, \mathrm{~b}, \mathrm{c}, \mathrm{~d}, \mathrm{e}, \mathrm{f}\} \\
& \mathrm{D}=\{\operatorname{SID} 01, \operatorname{SID} 02, \operatorname{SID} 03\} \\
& \text { SID01 }=\langle\mathrm{c}, \mathrm{f}, \mathrm{a}, \mathrm{~b}, \mathrm{e}, \mathrm{c}, \mathrm{~b}, \mathrm{a}, \mathrm{c}, \mathrm{~d}, \mathrm{a}\rangle \\
& \text { SID02 }=\langle\mathrm{f}, \mathrm{~b}, \mathrm{~d}, \mathrm{a}, \mathrm{c}, \mathrm{a}, \mathrm{c}, \mathrm{f}\rangle \\
& \text { SID03 }=\langle\mathrm{b}, \mathrm{e}, \mathrm{~d}, \mathrm{e}, \mathrm{f}, \mathrm{a}\rangle
\end{aligned}
$$

## III.1. Function Definition

The procedure of the AFiS algorithm started by defining the functions in a sequence. The function can be time series, procedures, or just function without any labels. The time series can contain data of days, months, or years, with value of weather, selling of goods, or others. The procedures can be a concept, such as the concept of $A-B-C-D$ in gamelan music implemented in this research, with notes as the value.

The number of function is determined based on each case. $F$ denotes a series of functions, and $T F$ denotes the total number of functions, then $F=\left(F_{1}, F_{2}, \ldots, F_{n}\right)$, and the total number of functions is the length of $F(T F=$ F.Length). For an example, define 3 functions in a sequence, then: $F=\left\{\mathrm{F}_{1}, \mathrm{~F}_{2}, \mathrm{~F}_{3}\right\}$, and $T F=3$.

## III.2. Data Partition

The second phase is data partition, where each sequence is partitioned with multiples of $T F$. The data partition for each sequence can be formulated as the pseudo-code of the data partition below:

```
S : sequence
TSI : total number of itemsets in a
    sequence
TF : total number of functions
P}\quad: data partition
n}=
While ( n < (TSI / TF ) ) {
        P [n] = [ ]
        n++
}
For ( n = 0; n < TSI; n++ ){
    For (k = 0; k < TF; k++){
        P [n] [k]=S [ (k*TF) + n]
    }
}
```

The value of $T F$ is 3 , then the SID01 containing 11 items will have 3 partitions; SID02 will have 3 partitions; SID03 will have 2 partitions, as shown in Fig. 2, with partition ID abbreviated as PID.

| SID | PID | P |
| :---: | :---: | :---: |
| 001 | 01 | $\langle\mathrm{c}, \mathrm{f}, \mathrm{a}\rangle$ |
|  | 02 | $\langle\mathrm{~b}, \mathrm{e}, \mathrm{c}\rangle$ |
|  | 03 | $\langle\mathrm{~b}, \mathrm{a}, \mathrm{c}\rangle$ |
|  | 04 | $\langle\mathrm{~d}, \mathrm{a}\rangle$ |
| 02 | 01 | $\langle\mathrm{f}, \mathrm{b}, \mathrm{d}\rangle$ |
|  | 02 | $\langle\mathrm{a}, \mathrm{c}, \mathrm{a}\rangle$ |
|  | 03 | $\langle\mathrm{c}, \mathrm{f}\rangle$ |
| 03 | 01 | $\langle\mathrm{~b}, \mathrm{e}, \mathrm{d}\rangle$ |
|  | 02 | $\langle\mathrm{e}, \mathrm{f}, \mathrm{a}\rangle$ |

Fig. 2. Result of data partition with $\mathrm{TF}=3$

## III.3. Sequential Pattern Creation

The sequential patterns are created by chaining functions. Each function is filled with itemsets from the sequences based on their order in the data partition. This procedure makes the first function is filled with itemsets which are in the first order of the partition, the second function is filled with itemsets which are in the second order of the partition, and so on.

There are 3 functions of $F 1, F 2, F 3$ and 3 sequences which have been partitioned. The result of itemsets defining of each function can be seen in Fig. 3.

| SID | PID | P |
| :---: | :---: | :---: |
| 01 | 01 | $<\mathrm{c}, \mathrm{f}, \mathrm{a}\rangle$ |
|  | 02 | $<\mathrm{b}, \mathrm{e}, \mathrm{c}\rangle$ |
|  | 03 | $<\mathrm{b}, \mathrm{a}, \mathrm{c}\rangle$ |
|  | 04 | $<\mathrm{d}, \mathrm{a}\rangle$ |
| 02 | 01 | $<\mathrm{f}, \mathrm{b}, \mathrm{d}\rangle$ |
|  | 02 | $\langle\mathrm{a}, \mathrm{c}, \mathrm{a}\rangle$ |
|  | 03 | $\langle\mathrm{c}, \mathrm{f}\rangle$ |
| 03 | 03 | $<\mathrm{b}, \mathrm{e}, \mathrm{d}\rangle$ |
|  | 01 | $<\mathrm{e}, \mathrm{f}, \mathrm{a}\rangle$ |


| F1 | F2 | F3 |
| :---: | :---: | :---: |
| c | f | a |
| b | e | c |
| b | $a$ | $c$ |
| d | $a$ | - |
| f | b | d |
| a | $c$ | $a$ |
| c | f | - |
| b | $e$ | $d$ |
| $e$ | f | $a$ |

Fig. 3. The result of defining the itemsets of each function
The process of defining itemsets of each function based on partition data can be formulated as the pseudocode below:

```
F : functions
P : data partition
TP : total number of partitions
    \(\mathrm{n}=0\)
    while ( \(\mathrm{n}<\mathrm{F}\).length) \(\{\)
        for \((k=0 ; k<T P ; k++)\{\)
            \(\mathrm{F}[\mathrm{n}][\mathrm{k}]=\mathrm{P}[\mathrm{k}][\mathrm{n}]\)
        \}
        n++
    \}
```

The sequential patterns are built by chaining functions. Since chaining functions requires at least 2 functions, then the sequential patterns started with 2-
itemsets up to ( $T F-1$ )-itemsets. There are 3 functions defined in this example, thus there are 2 sequential patterns containing 2 -itemsets, and 3 -itemsets.

The sequential pattern of 2-itemsets consists of $\langle\mathrm{F} 1$, F2>, <F2, F3>, <F3, F1*>, with asterisk symbols standing for the next partition. The sequential pattern of 3-itemsets consists of <F1, F2, F3>, <F2, F3, F1*>, <F3, $\mathrm{F} 1^{*}, \mathrm{~F} 2^{*}>$. Fig. 4 shows the itemsets of the sequential patterns of 2-itemsets, and 3-itemsets.

| Itemsets of Sequential Pattern of 2-Itemsets |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| SID | PID | $\begin{aligned} & <\mathbf{F 1}, \\ & \text { F2 }> \end{aligned}$ | $\begin{aligned} & <\mathbf{F} 2, \\ & \mathbf{F 3}> \end{aligned}$ | $\begin{aligned} & <\mathbf{F 3}, \\ & \text { F1*> } \end{aligned}$ |
| 01 | 01 | <c, f> | <f, a> | <a, b> |
|  | 02 | <b, e> | $\langle\mathrm{e}, \mathrm{c}\rangle$ | $<\mathrm{c}, \mathrm{b}>$ |
|  | 03 | <b, a> | $<\mathrm{a}, \mathrm{c}>$ | $\langle\mathrm{c}, \mathrm{d}\rangle$ |
|  | 04 | <d, a> | <a> | - |
| 02 | 01 | <f, b> | $<\mathrm{b}, \mathrm{d}\rangle$ | <d, a $>$ |
|  | 02 | $\langle\mathrm{a}, \mathrm{c}\rangle$ | <c, a> | <a, c> |
|  | 03 | $<\mathrm{c}, \mathrm{f}>$ | <f $>$ | - |
| 03 | 01 | <b, e> | $<\mathrm{e}, \mathrm{d}\rangle$ | <d, e> |
|  | 02 | <e, f> | <f, a> | <a> |


| SID | PID | $\begin{gathered} \langle\mathbf{F 1}, \mathbf{F} 2, \\ \text { F3 }> \end{gathered}$ | $\begin{gathered} \langle\mathbf{F} 2, ~ F 3, \\ \mathbf{F 1}^{*}> \end{gathered}$ | $\begin{gathered} <\mathbf{F 3}, \mathbf{F 1} 1^{*}, \\ \text { F2* } \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: |
| 01 | 01 | <c, f, a> | <f, a, b> | <a, b, e> |
|  | 02 | $<\mathrm{b}, \mathrm{e}, \mathrm{c}\rangle$ | $<\mathrm{e}, \mathrm{c}, \mathrm{b}$ > | $<\mathrm{c}, \mathrm{b}, \mathrm{a}>$ |
|  | 03 | $\langle\mathrm{b}, \mathrm{a}, \mathrm{c}\rangle$ | $\langle\mathrm{a}, \mathrm{c}, \mathrm{d}\rangle$ | <c, d, a> |
|  | 04 | $<\mathrm{d}, \mathrm{a}>$ | $<\mathrm{a}>$ | - |
| 02 | 01 | $\langle\mathrm{f}, \mathrm{b}, \mathrm{d}\rangle$ | $<\mathrm{b}, \mathrm{d}, \mathrm{a}\rangle$ | <d, a, c> |
|  | 02 | <a, c, a $>$ | $\langle\mathrm{c}, \mathrm{a}, \mathrm{c}\rangle$ | $<\mathrm{a}, \mathrm{c}, \mathrm{f}>$ |
|  | 03 | <c, f> | <f $>$ | - |
| 03 | 01 | $<\mathrm{b}, \mathrm{e}, \mathrm{d}\rangle$ | <e, d, e> | <d, e, f> |
|  | 02 | <e, f, a> | <f, a> | <a> |

Fig. 4. Sequential patterns created by chaining the functions
The next phase is candidate selection, where the itemsets with a length not equal to the length of the sequential pattern (k-itemsets) is eliminated from the list. For example, itemsets $\langle\mathrm{a}>$ in sequential pattern of 2itemsets <F2, F3> is eliminated from the list. Fig. 5 shows the results of candidate selection.

| Candidates of 2-Itemsets |  |  |  |
| :---: | :---: | :---: | :---: |
| SID | $\begin{aligned} & <\text { F1, } \\ & \text { F2> } \end{aligned}$ | $\begin{aligned} & <\mathrm{F} 2, \\ & \text { F3> } \end{aligned}$ | $\begin{aligned} & <\text { F3, } \\ & \text { F1*> } \end{aligned}$ |
| 01 | <c, f> | <f, a> | <a, b> |
|  | $\langle\mathrm{b}, \mathrm{e}\rangle$ | $\langle\mathrm{e}, \mathrm{c}\rangle$ | $<\mathrm{c}, \mathrm{b}\rangle$ |
|  | < b, a> | $\langle\mathrm{a}, \mathrm{c}\rangle$ | $\langle\mathrm{c}, \mathrm{d}\rangle$ |
|  | <d, a> | - | - |
| 02 | <f, b> | < $\mathrm{b}, \mathrm{d}\rangle$ | <d, a> |
|  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | <c, a> | $\langle\mathrm{a}, \mathrm{c}\rangle$ |
|  | $<\mathrm{c}, \mathrm{f}>$ | - | - |
| 03 | $<\mathrm{b}, \mathrm{e}\rangle$ | <e, d> | <d, e> |
|  | $<\mathrm{e}, \mathrm{f}>$ | <f, a> | - |


| Candidates of 3-Itemsets |  |  |  |
| :---: | :---: | :---: | :---: |
| SID | $\begin{gathered} \langle\text { F1, F2, }, \\ \text { F3 }> \end{gathered}$ | $\begin{gathered} \langle\text { F2, F3, } \\ \text { F1*> } \end{gathered}$ | $\begin{gathered} \langle\text { F3, F1* } \\ \text { F2*> } \end{gathered}$ |
| 01 | <c, f, a> | <f, a, b> | <a, b, e> |
|  | $\langle\mathrm{b}, \mathrm{e}, \mathrm{c}\rangle$ | $<\mathrm{e}, \mathrm{c}, \mathrm{b}\rangle$ | <c, b, a> |
|  | $\langle\mathrm{b}, \mathrm{a}, \mathrm{c}\rangle$ | $\langle\mathrm{a}, \mathrm{c}, \mathrm{d}\rangle$ | <c, d, a> |
| 02 | $\langle\mathrm{f}, \mathrm{b}, \mathrm{d}\rangle$ | <b, d, a> | <d, a, c> |
|  | <a, c, a> | <c, a, c> | $<\mathrm{a}, \mathrm{c}, \mathrm{f}>$ |
| 03 | $\langle\mathrm{b}, \mathrm{e}, \mathrm{d}\rangle$ | <e, d, e> | <d, e, f $>$ |
|  | <e, f, a> | - | - |

Fig. 5. Candidates of sequential patterns

## III.4. Support Counting

The sequence defined as frequent is measured using minimum support value. An itemset must have at least 1 transaction to be defined as frequent, if the given minimum support is 1 . The pseudocode below is used to find frequent sequence:

```
LN : frequent sequences
C : sequent itemsets
minsup : minimum support
for (k=0; k < TC; k++ ){
    if( C [k] >= minsup ){
        LN.push ( s [k])
    }
}
```

For an example, if the given minimum support is 1 , then all candidates are frequent.

After counting the support, the next step is defining the dominant predicate for the itemset by measuring its weight. The formula below is used to measure the weight of the itemset, where $W$ stands for weight, $T I$ stands for the total number of itemsets for each pattern, and TC stands for the total number of candidates:

$$
\begin{equation*}
W=\frac{T I}{T C} \tag{1}
\end{equation*}
$$

Fig. 6 and Fig. 7 show the weight of sequence of 2itemset and 3 -itemsets.
Weight of 2-Itemsets $\langle$ F1, F2 $\rangle$

| SID |  |  |  |  |  | I | TI | TC | W |
| :--- | :--- | ---: | ---: | ---: | :---: | :---: | :---: | :---: | :---: |
| 01 | $\langle\mathrm{c}, \mathrm{f}\rangle$ | 1 | 4 | 0.25 |  |  |  |  |  |
|  | $<\mathrm{b}, \mathrm{e}\rangle$ | 1 | 4 | 0.25 |  |  |  |  |  |
|  | $<\mathrm{b}, \mathrm{a}\rangle$ | 1 | 4 | 0.25 |  |  |  |  |  |
|  | $\langle\mathrm{~d}, \mathrm{a}\rangle$ | 1 | 4 | 0.25 |  |  |  |  |  |
| 02 | $\langle\mathrm{f}, \mathrm{b}\rangle$ | 1 | 3 | 0.33 |  |  |  |  |  |
|  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 1 | 3 | 0.33 |  |  |  |  |  |
|  | $\langle\mathrm{c}, \mathrm{f}\rangle$ | 1 | 3 | 0.33 |  |  |  |  |  |
| 03 | $<\mathrm{b}, \mathrm{e}\rangle$ | 1 | 2 | 0.5 |  |  |  |  |  |
|  | $<\mathrm{e}, \mathrm{f}\rangle$ | 1 | 2 | 0.5 |  |  |  |  |  |


| Weight of 2-Itemsets < F2, F3> |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| SID | I | TI | TC | W |
| 01 | <f, a> | 1 | 3 | 0.33 |
|  | <e, c> | 1 | 3 | 0.33 |
|  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 1 | 3 | 0.33 |
| 02 | $\langle b, d>$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{c}, \mathrm{a}\rangle$ | 1 | 2 | 0.5 |
| 03 | <e, d> | 1 | 2 | 0.5 |
|  | $\langle f, a\rangle$ | 1 | 2 | 0.5 |

Weight of 2-Itemsets $\left\langle\mathbf{F 3}, \mathbf{F 1} \mathbf{F}^{*}\right\rangle$

| SID | I | TI | TC | W |
| :--- | :--- | ---: | ---: | ---: |
| 01 | $\langle\mathrm{a}, \mathrm{b}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{c}, \mathrm{b}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{c}, \mathrm{d}\rangle$ | 1 | 3 | 0.33 |
| 02 | $\langle\mathrm{~d}, \mathrm{a}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{~d}, \mathrm{e}\rangle$ | 1 | 1 | 1 |

Fig. 6. Weight of the sequences of 2-itemsets
Weight of 3-Itemsets $\langle$ F1, F2, F3 $\rangle$

| SID | I | TI | TC | W |
| :--- | :--- | ---: | ---: | ---: |
| 01 | $\langle\mathrm{c}, \mathrm{f}, \mathrm{a}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{~b}, \mathrm{e}, \mathrm{c}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{~b}, \mathrm{a}, \mathrm{c}\rangle$ | 1 | 3 | 0.33 |
| 02 | $\langle\mathrm{f}, \mathrm{b}, \mathrm{d}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{a}, \mathrm{c}, \mathrm{a}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{~b}, \mathrm{e}, \mathrm{d}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{e}, \mathrm{f}, \mathrm{a}\rangle$ | 1 | 2 | 0.5 |

Weight of 3-Itemsets $\left\langle\mathbf{F 2}, \mathbf{F 3}, \mathbf{F 1}{ }^{*}\right\rangle$

| SID | I | TI | TC | W |
| :--- | :---: | ---: | :--- | :--- |
|  | $\langle\mathrm{f}, \mathrm{a}, \mathrm{b}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{e}, \mathrm{c}, \mathrm{b}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{a}, \mathrm{c}, \mathrm{d}\rangle$ | 1 | 3 | 0.33 |
| 2 | $\langle\mathrm{~b}, \mathrm{~d}, \mathrm{a}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{c}, \mathrm{a}, \mathrm{c}\rangle$ | 1 | 2 | 0.5 |
| 03 | $\langle\mathrm{e}, \mathrm{d}, \mathrm{e}\rangle$ | 1 | 1 | 1 |

Weight of 3-Itemsets $\left\langle\mathbf{F 3}, \mathbf{F 1 *}\right.$ * $\left.\mathbf{F 2}^{*}\right\rangle$

| SID | I | TI | TC | W |
| :--- | :--- | ---: | ---: | ---: |
| 01 | $\langle\mathrm{a}, \mathrm{b}, \mathrm{e}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{c}, \mathrm{b}, \mathrm{a}\rangle$ | 1 | 3 | 0.33 |
|  | $\langle\mathrm{c}, \mathrm{d}, \mathrm{a}\rangle$ | 1 | 3 | 0.33 |
| 02 | $\langle\mathrm{~d}, \mathrm{a}, \mathrm{c}\rangle$ | 1 | 2 | 0.5 |
|  | $\langle\mathrm{a}, \mathrm{c}, \mathrm{f}\rangle$ | 1 | 2 | 0.5 |
| 03 | $\langle\mathrm{~d}, \mathrm{e}, \mathrm{f}\rangle$ | 1 | 1 | 1 |

Fig. 7. Weight of the sequences of 3 -itemsets
The next step is concatenating the frequent sequences to measure their total weight. For the itemsets containing the same items, their weights are summed up. Then, the weight of each itemset is divided into the total number of sequences. The formula below is used to measure the total weight of the itemset, where $T W$ stands for total weight, $W$ stands for the weight of itemsets, and $T S$ stands for the total number of sequences:

$$
\begin{equation*}
T W=\frac{W}{T S} \tag{2}
\end{equation*}
$$

Fig. 8 shows the total weight of the sequences of 2itemsets and 3-itemsets.

| Total weight of 2-Itemsets$\langle\mathrm{F} 1, \mathrm{~F} 2\rangle$ |  |  |  |
| :---: | :---: | :---: | :---: |
| I | W | SW | TW |
| <a, c> | 0.33 | 0.33 | 0.110 |
| $\langle\mathrm{b}, \mathrm{a}$ > | 0.25 | 0.25 | 0.083 |
| <b, e> | 0.5 | 0.75 | 0.250 |
|  | 0.25 |  |  |
| $<\mathrm{c}, \mathrm{f}>$ | 0.33 | 0.58 | 0.193 |
|  | 0.25 |  |  |
| <d, a> | 0.25 | 0.25 | 0.083 |
| <e, f $>$ | 0.5 | 0.5 | 0.167 |
| <f, b> | 0.33 | 0.33 | 0.110 |

Total Weight of 3-Itemsets

| $\langle$ F1, F2, F3 $\rangle$ |  |  |  |
| :---: | :--- | :--- | :--- |
| I | W | SW | TW |
| $\langle\mathrm{a}, \mathrm{c}, \mathrm{a}\rangle$ | 0.5 | 0.5 | 0.167 |
| $\langle\mathrm{~b}, \mathrm{a}, \mathrm{c}\rangle$ | 0.33 | 0.33 | 0.110 |
| $\langle\mathrm{~b}, \mathrm{e}, \mathrm{c}\rangle$ | 0.33 | 0.33 | 0.110 |
| $\langle\mathrm{~b}, \mathrm{e}, \mathrm{d}\rangle$ | 0.5 | 0.5 | 0.167 |
| $\langle\mathrm{c}, \mathrm{f}, \mathrm{a}\rangle$ | 0.33 | 0.33 | 0.110 |
| $\langle\mathrm{e}, \mathrm{f}, \mathrm{a}\rangle$ | 0.5 | 0.5 | 0.167 |
| $\langle\mathrm{f}, \mathrm{b}, \mathrm{d}\rangle$ | 0.5 | 0.5 | 0.167 |

Total Weight of 3-Itemsets


| $\langle\mathbf{F 2 , F 3}, \mathbf{F 1 *}\rangle$ |  |  |  |  |
| :--- | :--- | :--- | :--- | :---: |
| I | $\mathbf{W}$ | SW | TW |  |
| $\langle\mathrm{a}, \mathrm{c}, \mathrm{d}\rangle$ | 0.33 | 0.33 | 0.110 |  |
| $\langle\mathrm{~b}, \mathrm{~d}, \mathrm{a}\rangle$ | 0.5 | 0.5 | 0.167 |  |
| $\langle\mathrm{c}, \mathrm{a}, \mathrm{c}\rangle$ | 0.5 | 0.5 | 0.167 |  |
| $\langle\mathrm{e}, \mathrm{c}, \mathrm{b}\rangle$ | 0.33 | 0.33 | 0.110 |  |
| $\langle\mathrm{e}, \mathrm{d}, \mathrm{e}\rangle$ | 1 | 1 | 0.333 |  |
| $\langle\mathrm{f}, \mathrm{a}, \mathrm{b}\rangle$ | 0.33 | 0.33 | 0.110 |  |

Total Weight of 3-Itemsets


| $\left\langle\mathbf{F 3}, \mathbf{F 1}^{*}, \mathbf{F 2}^{*}\right\rangle$ |  |  |  |  |
| :--- | :--- | :--- | :--- | :---: |
| $\mathbf{I}$ | $\mathbf{W}$ | SW | TW |  |
| $\langle\mathrm{a}, \mathrm{b}, \mathrm{e}\rangle$ | 0.33 | 0.33 | 0.110 |  |
| $\langle\mathrm{a}, \mathrm{c}, \mathrm{f}\rangle$ | 0.5 | 0.5 | 0.167 |  |
| $\langle\mathrm{c}, \mathrm{b}, \mathrm{a}\rangle$ | 0.33 | 0.33 | 0.110 |  |
| $\langle\mathrm{c}, \mathrm{d}, \mathrm{a}\rangle$ | 0.33 | 0.33 | 0.110 |  |
| $\langle\mathrm{~d}, \mathrm{a}, \mathrm{c}\rangle$ | 0.5 | 0.5 | 0.167 |  |
| $\langle\mathrm{~d}, \mathrm{e}, \mathrm{f}\rangle$ | 1 | 1 | 0.333 |  |

Fig. 8. Total weight of the 2 -itemsets and 3 -itemsets

## III.5. Prune Phase

Prune phase is used to chain itemsets based on functions. For example, there is an itemset <a, c> in the function $\langle\mathrm{F} 1, \mathrm{~F} 2\rangle$, so the itemset in the function $<\mathrm{F} 2$, F3> must start with itemset $c$, such as <c, a>, and the itemset in the function $\langle\mathrm{F} 3, \mathrm{~F} 1 *\rangle$ must start with item a, such as <a, b>, and <a, c>. The last itemset in the itemset of $\left\langle\mathrm{F} 3, \mathrm{~F} 1^{*}\right\rangle$ is used as a reference to set the first item in the <F1, F2> itemset. Fig. 9 and Fig. 10 show the results of the prune phase.

| F1 | <F1, F2> |  |  |  |  |  | F1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | <F2, F3> |  | $<$ F3, F1*> |  |  |
|  | I | TW | I | TW | I | TW |  |
| a | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 0.110 | $\langle\mathrm{c}, \mathrm{a}\rangle$ | 0.167 | <a, b> | 0.110 | b |
|  |  |  |  |  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 0.167 | c |
| b | $<\mathrm{b}, \mathrm{a}>$ | 0.083 | <a, c> | 0.110 | <c, b> | 0.110 | b |
|  |  |  |  |  | $\langle\mathrm{c}, \mathrm{d}\rangle$ | 0.110 | d |
|  | $\langle\mathrm{b}, \mathrm{e}>$ | 0.250 | $\langle\mathrm{e}, \mathrm{c}\rangle$ | 0.110 | $\langle\mathrm{c}, \mathrm{b}\rangle$ | 0.110 | b |
|  |  |  |  |  | <c, d> | 0.110 | d |
|  |  |  | $\langle\mathrm{e}, \mathrm{d}\rangle$ | 0.167 | <d, a> | 0.167 | a |
|  |  |  |  |  | <d, e> | 0.333 | e |
| c | $<\mathrm{c}, \mathrm{f}>$ | 0.193 | <f, a> | 0.277 | <a, b> | 0.110 | b |
|  |  |  |  |  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 0.167 | c |
| d | <d, a> | 0.083 | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 0.110 | <c, b> | 0.110 | b |
|  |  |  |  |  | $\langle\mathrm{c}, \mathrm{d}\rangle$ | 0.110 | d |
| e | <e, f> | 0.167 | <f, a> | 0.277 | $\langle\mathrm{a}, \mathrm{b}\rangle$ | 0.110 | b |
|  |  |  |  |  | $\langle\mathrm{a}, \mathrm{c}\rangle$ | 0.167 | c |
| f | <f, b> | 0.110 | $<\mathrm{b}, \mathrm{d}>$ | 0.167 | <d, a> | 0.167 | a |
|  |  |  |  |  | <d, e> | 0.333 | e |

Fig. 9. Prunes of 2-itemsets

| Prunes of 3-Itemsets |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \langle\mathrm{F} 1, \\ & \text { F2> } \end{aligned}$ | <F1, F2, F3> |  | <F2, F3, F1*> |  | <F3, F1*, F2* ${ }^{\text {* }}$ |  | $\begin{aligned} & \langle\text { F1, } \\ & \text { F2> } \end{aligned}$ |
|  | I | TW | I | TW | I | TW |  |
| ac | <a, c, a> | 0.167 | <c, a, c> | 0.167 | <a, c, f $>$ | 0.167 | cf |
| ba | $\langle\mathrm{b}, \mathrm{a}, \mathrm{c}\rangle$ | 0.110 | <a, c, d> | 0.110 | <c, d, a> | 0.110 | da |
| be | $\langle\mathrm{b}, \mathrm{e}, \mathrm{c}\rangle$ | 0.110 | <e, c, b> | 0.110 | <c, b, a> | 0.110 | ba |
|  | $\langle\mathrm{b}, \mathrm{e}, \mathrm{d}\rangle$ | 0.167 | <e, d, e> | 0.333 | <d, e, f $>$ | 0.333 | ef |
| cf | <c, f, a> | 0.110 | <f, a, b> | 0.110 | $\langle\mathrm{a}, \mathrm{b}, \mathrm{e}\rangle$ | 0.110 | be |
| ef | $\langle\mathrm{e}, \mathrm{f}, \mathrm{a}\rangle$ | 0.167 |  |  |  |  |  |
| fb | $\langle f, b, d\rangle$ | 0.167 | $\langle\mathrm{b}, \mathrm{d}, \mathrm{a}>$ | 0.167 | <d, a, c> | 0.167 | ac |

Fig. 10. Prunes of 3-itemsets

## III.6. Production Rules

The rules are set based on the prunes.For example, the itemset <a, c> in the function <F1, F2> chains to the itemset $\langle\mathrm{c}$, $\mathrm{a}>$ in the function $\langle\mathrm{F} 2, \mathrm{~F} 3$ >, and it continues to chain to itemsets $\langle\mathrm{a}, \mathrm{b}\rangle$, and $\langle\mathrm{a}, \mathrm{c}\rangle$ in the function <F3, F1*>. The itemset 〈b> in 〈a, b> of the function $<\mathrm{F} 3, \mathrm{~F} 1^{*}>$ is then chained to itemsets <b, a>, <b, e> of the function <F1, F2>. This is also done to the itemset $\langle\mathrm{c}\rangle$ in <a, c> of the function <F3, F1*>, where it can be chained to the itemset $\langle\mathrm{c}, \mathrm{f}\rangle$ of the function $<\mathrm{F} 1$, F2>. Below are the examples of the 2 -itemsets rules set based on the prunes:

```
Start
    IF F1 is a
    THEN F2 is }
    IF F1 is }
    THEN F2 is a OR e
```

```
    IF F1 is d
    THEN F2 is }
    IF F1 is }
    THEN F2 is }
    IF F1 is f
    THEN F2 is }
    ...
    IF F2 is }
    THEN F3 is cOR }
    IF F2 is f
    THEN F3 is }
    IF F3 is a
    THEN F1* is bOR }
    ...
    IF F3 is d
    THEN F1* is }d\mathrm{ OR }
End
```

The next examples are the 3-itemsets rules set based on the prunes.

```
Start
    IF F1 is }
    THEN F2 is }
    ANDF3 is }
```

    IF F1 is \(b\) AND F2 is \(a\)
    THEN F3 is \(c\)
    IF F1 is \(b\) AND F2 is \(e\)
    THEN F3 is \(c \mathbf{O R} d\)
    \(\cdots\)
    IF F1 is \(f\) AND F2 is \(b\)
    THEN F3 is \(d\)
    IF F3 is \(a\) AND F1* is \(c\)
    THEN F2* is \(f\)
    IF F3 is \(d\) AND F1* is \(a\)
    THEN F2* is \(c\)
    IF F3 is \(d\) AND F1* is \(e\)
    THEN F2* is \(f\)
    
## End

## IV. Simulation Work and Analysis

AFiS algorithm is implemented to build melodic feature knowledge of gamelan music, and to set rules of note sequences. The dataset is limited to gending ladrang
laras slendro pathet nem. 15 gendings entitled Konda, Rangsang, Plupuh, Erang-Erang, Jong Layar, Mangu, Peksi Kuwung, Sobah, Medang Miring, Lung Gadung Pel, Kembang Gadhung Ngayangan, Alas Kobong, Binar, Kandha,Liwung were collected as a dataset. The discussion below describes the implementation of AFiS algorithm procedures in building melodic feature knowledge, and setting rules of note sequences in gamelan music.

Most of the gendings collected to be used as dataset have gatras with balungan nibani types. According to the characteristics of balungan nibani, most gatras use pins (dot notes) in their note sequences. For example, a gending entitled Medang Miring seen in Fig. 11, shows that pins are used in most gatras, and only 2 gatras which have type of balungan mlaku (without pins in their note sequences).


Fig. 11. Pins in note sequences
The analysis is to find the causality correlation among note sequences, and pins are not included as note sequences. The type of balungan nibani inserting pins into a gatra cannot be used as a dataset. Data cleansing is conducted by converting balungan nibani into balungan mlaku. Pins are removed from note sequences, and new gatras are constructed through a concatenating order.The process of data cleansing is implemented to all 15 gendings used as dataset. Fig. 12 shows the newgatras used as dataset in gending entitled 'Medang Miring'.

```
Ladrang 'Medang Miring'
Laras Slendro Pathet Nem
3235 3235
1216 2161
6535
```

Fig. 12. New gatras from conversion process
As the first phase of the procedure of AFiS algorithm, after the data cleansing process, the functions in a note sequence are defined. The note sequence in a gending is arranged based on gatra. Gatra is the smallest unit consisting of 4 sabetan balungan (beats). Fig. 13 shows an illustration of gatra, with sabetan balungan symbolized as $A, B, C, D$.


Fig. 13. Illustration of gatra

The concept of gatra is known as the $A-B-C-D$ concept: $A$ represents maju (forward), $B$ represents mundur (back), $C$ represents maju (forward), and $D$ represents seleh (terminal/the end point of a journey). This concept implies the existence of a hierarchy of functions of each part of the gatra, where $D$ is the musical point reference or the strongest part, $B$ is the second strongest part, $A$ is the third strongest part, and $C$ is the weakest part.

The concept of $A-B-C-D$ in a gatra is used as a reference in constructing the functions, where the first note represents function $A$, the second note represents function $B$, the third note represents function $C$, and the fourth note represents function $D$.

Data partition is used to adjust the notes to the functions in sequences. There are 4 functions used in this analysis. Note sequences data are partitioned with length at 4 , as the length of the number of functions as a representation of the concept of $A-B-C-D$. After conducing data partition using formula 1 , each partition contains 4 notes, where each note fills the value of the function based on its sequence to create sequential patterns.

Data partition and sequential patterns creation are implemented to all 15 gendings in the dataset. Table I shows the example of data partition and sequential patterns creation of a gending entitled Konda.

TABLE I
Example Of Data Partition And Sequential Patterns Creation

| NO | $\begin{gathered} \text { DATA } \\ \text { PARTITION } \end{gathered}$ | A | B | C | D |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | <3, 5, 3, 2> | 3 | 5 | 3 | 2 |
| 2 | <6, 1, 3, 2> | 6 | 1 | 3 | 2 |
| 3 | <3, 2, 3, 2> | 3 | 2 | 3 | 2 |
| 4 | <6, 5, 3, 2> | 6 | 5 | 3 | 2 |
| 5 | <3, 5, 3, 2> | 3 | 5 | 3 | 2 |
| 6 | <3, 2, 1, 6> | 3 | 2 | 1 | 6 |
| 7 | <3, 2, 1, 6> | 3 | 2 | 1 | 6 |
| 8 | <3, 5, 3, 2> | 3 | 5 | 3 | 2 |
| 9 | <3, 5, 3, 2> | 3 | 5 | 3 | 2 |
| 10 | <5, 6, 3, 5> | 5 | 6 | 3 | 5 |
| 11 | $<6,5,6,5\rangle$ | 6 | 5 | 6 | 5 |
| 12 | <1, 6, 3, 2> | 1 | 6 | 3 | 2 |
| 13 | <3, 5, 3, 2> | 3 | 5 | 3 | 2 |
| 14 | <5, 3, 1, 6> | 5 | 3 | 1 | 6 |
| 15 | <3, 2, 1, 6> | 3 | 2 | 1 | 6 |
| 16 | <3, 5, 3, 2> | 3 | 5 | 3 | 2 |

The notes sequences data which have been partitioned based on the $A-B-C-D$ function, and defined as sequential patterns, are then constructed to chain functions.

The sequential patterns started with 2 -itemsets up to 4itemsets. The sequential pattern of 2 -itemsets consists of function chaining of $\langle\mathrm{A}, \mathrm{B}\rangle,\langle\mathrm{B}, \mathrm{C}\rangle,\langle\mathrm{C}, \mathrm{D}\rangle,\left\langle\mathrm{D}, \mathrm{A}^{*}\right\rangle$, with asterisk symbols that stand for the next partition. The sequential pattern of 3 -itemsets consists of function chaining of $\langle\mathrm{A}, \mathrm{B}, \mathrm{C}\rangle,\langle\mathrm{B}, \mathrm{C}, \mathrm{D}\rangle,\left\langle\mathrm{C}, \mathrm{D}, \mathrm{A}^{*}\right\rangle,\left\langle\mathrm{D}, \mathrm{A}^{*}\right.$, B*).

The sequential pattern of 4-itemsets consists of function chaining of $\langle\mathrm{A}, \mathrm{B}, \mathrm{C}, \mathrm{D}\rangle,\langle\mathrm{B}, \mathrm{C}, \mathrm{D}, \mathrm{A}\rangle,\langle\mathrm{C}, \mathrm{D}$, $\left.A^{*}, B^{*}\right\rangle,\left\langle D, A^{*}, B^{*}, C^{*}\right\rangle$.

The process of creating sequential patterns is implemented to all 15 gendings in the dataset.

The following tables show examples of sequential patterns of 2-itemsets, 3-itemsets, and 4-itemsets of a gending entitled Konda.

TABLE II
Example Of Sequential Pattern Of 2-Itemsets Of A Gending Entitled Konda

| <A, B> | <B, C> | <C, D> | <D, A*> |
| :---: | :---: | :---: | :---: |
| <3, 5> | <5, 3> | <3, 2> | <2, 6> |
| <6, 1> | <1, 3> | <3, 2> | <2, 3> |
| <3, 2> | <2, 3> | <3, 2> | <2, 6> |
| <6, 5> | $\langle 5,3\rangle$ | <3, 2> | <2, 3> |
| <3, 5> | <5, 3> | <3, 2> | <2, 3> |
| < 3 , 5 > |  |  | <2> |

TABLE III
Example Of Sequential Patterns Of 3-Itemsets Of A Gending
Entitled Konda

| <A, B, C> | <B, C, D> | <C, D, A*> | <D, A*, B*> |
| :---: | :---: | :---: | :---: |
| $\langle 3,5,3\rangle$ | <5, 3, 2> | <3, 2, 6> | <2, 6, 1> |
| <6, 1, 3> | $\langle 1,3,2\rangle$ | <3, 2, 3> | <2, 3, 2> |
| <3, 2, 3> | <2, 3, 2> | <3, 2, 6> | <2, 6, 5> |
| $<6,5,3>$ | <5, 3, 2> | <3, 2, 3> | <2, 3, 5> |
| <3, 5, 3> | <5, 3, 2> | <3, 2, 3> | <2, 3, 2> |
| < $\ldots$. 5,3$\rangle$ | $\ldots$ $\langle 5,3,2>$ | . $\langle 3,2>$ | <2> |

TABLE IV
Example Of Sequential Patterns Of 4-Itemsets Of A Gending Entitled Konda

| $\langle\mathbf{A}, \mathbf{B}, \mathbf{C}$, | $\langle\mathbf{B}, \mathbf{C}, \mathbf{D}$, | $\left\langle\mathbf{C}, \mathbf{D}, \mathbf{A}^{*}\right.$, | $\left\langle\mathbf{D}, \mathbf{A}^{*}, \mathbf{B}^{*}\right.$, |
| :---: | :---: | :---: | :---: |
| $\mathbf{D}\rangle$ | $\left.\mathbf{A}^{*}\right\rangle$ | $\left.\mathbf{B}^{*}\right\rangle$ | $\left.\mathbf{C}^{*}\right\rangle$ |
| $\langle 3,5,3,2\rangle$ | $\langle 5,3,2,6\rangle$ | $\langle 3,2,6,1\rangle$ | $\langle 2,6,1,3\rangle$ |
| $\langle 6,1,3,2\rangle$ | $\langle 1,3,2,3\rangle$ | $\langle 3,2,3,2\rangle$ | $\langle 2,3,2,3\rangle$ |
| $\langle 3,2,3,2\rangle$ | $\langle 2,3,2,6\rangle$ | $\langle 3,2,6,5\rangle$ | $\langle 2,6,5,3\rangle$ |
| $\ldots \ldots, 2\rangle$ | $\langle 5,3,2\rangle$ | $\langle 3,2\rangle$ | $\ldots$ |
| $\langle 3,5,3,2\rangle$ |  | $<2\rangle$ |  |

The next phase is candidate selection, where the itemsets with a length not equal to the length of its sequential pattern are eliminated.

For example, in the SID 01, the itemsets <2> in 2itemsets $\left\langle\mathrm{D}, \mathrm{A}^{*}\right\rangle,\langle 3,2\rangle$ in 3 -itemsets $\left\langle\mathrm{C}, \mathrm{D}, \mathrm{A}^{*}\right\rangle,\langle 2\rangle$ in the 3-itemsets $\left\langle\mathrm{D}, \mathrm{A}^{*}, \mathrm{~B}^{*}\right\rangle,\langle 5,3,2\rangle$ in the 4-itemsets $\left.<\mathrm{B}, \mathrm{C}, \mathrm{D}, \mathrm{A}^{*}\right\rangle,\langle 3,2\rangle$ in the 4 -itemsets <C, $\left.\mathrm{D}, \mathrm{A}^{*}, \mathrm{~B}^{*}\right\rangle$, and $\langle 2\rangle$ in the 4 -itemsets $\left\langle\mathrm{D}, \mathrm{A}^{*}, \mathrm{~B}^{*}, \mathrm{C}^{*}\right\rangle$, are eliminated.

The candidate selection process is implemented to all 15 gendings in the dataset.

The itemsets in all functions of each candidate are then measured using minimum support value. The minimum support is defined at 1 .

The results show that all the candidates of all 15 gendings are defined as frequent sequences, and then the weight of the frequent itemsets in all functions is measured using formula (1).

The concatenation process is then implemented to the itemsets in all 15 gendings based on each function. Then, the concatenation itemsets are measured for their total weight using formula (2).

Tables below show the frequent sequences of 2itemsets, 3 -itemsets, and 4 -itemsets resulted from concatenation process of 15 gendings.

TABLE V
Frequent Sequences Of 2-Itemsets Resulted From CONCATENATION PROCESS

|  | CONA |
| :---: | :--- |
| $\langle\mathbf{A}, \mathbf{B}\rangle$ | $\langle 1,2\rangle,\langle 1,6\rangle,\langle 2,1\rangle,\langle 2,3\rangle,\langle 2,6\rangle,\langle 3,1\rangle,\langle 3,2\rangle$, |
| $\langle\mathbf{B}, \mathbf{C}\rangle$ | $\langle 3,5\rangle,\langle 3,6\rangle,\langle 5,3\rangle,\langle 5,6\rangle,\langle 6,1\rangle,\langle 6,3\rangle,\langle 6,5\rangle$ |
|  | $\langle 1,2\rangle,\langle 1,3\rangle,\langle 1,5\rangle,\langle 1,6\rangle,\langle 2,1\rangle,\langle 2,3\rangle,\langle 2,5\rangle$, |
| $\langle\mathbf{C}, \mathbf{D}\rangle$ | $\langle 2,6\rangle,\langle 3,1\rangle,\langle 3,5\rangle,\langle 3,6\rangle,\langle 5,1\rangle, \ldots,\langle 6,5\rangle$ |
|  | $\langle 1,6\rangle,\langle 2,1\rangle,\langle 2,6\rangle,\langle 3,2\rangle,\langle 3,5\rangle,\langle 5,2\rangle,\langle 5,3\rangle$, |
| $\left\langle\mathbf{D}, \mathbf{A}^{*}\right\rangle$ | $\langle 5,6\rangle,\langle 6,1\rangle,\langle 6,5\rangle$ |
|  | $\langle 1,2\rangle,\langle 1,6\rangle,\langle 2,1\rangle,\langle 2,3\rangle,\langle 2,5\rangle,\langle 2,6\rangle,\langle 3,1\rangle$, |
| $\langle 3,2\rangle,\langle 3,5\rangle,\langle 5,1\rangle,\langle 5,2\rangle,\langle 5,3\rangle, \ldots,\langle 6,5\rangle$ |  |

TABLE VI
Frequent Sequences Of 3-Itemsetsresulted From CONCATENATION PROCESS

| Concatenation Process |  |
| :---: | :---: |
| <A, B, C> | $\begin{aligned} & \langle 1,2,1\rangle,\langle 1,6,3\rangle,\langle 1,6,5\rangle,\langle 2,1,2\rangle,\langle 2,1,5\rangle, \\ & \langle 2,1,6\rangle,\langle 2,3,1\rangle,\langle 2,3,6\rangle,\langle 2,6,5\rangle,\langle 3,1,3\rangle, \\ & \langle 3,1,5\rangle,\langle 3,2,1\rangle,\langle 3,2,3\rangle, \ldots,\langle 6,5,6\rangle \end{aligned}$ |
| <B, C, D> | $\begin{aligned} & \langle 1,2,6\rangle,\langle 1,3,2\rangle,\langle 1,5,3\rangle,\langle 1,5,6\rangle,\langle 1,6,1\rangle, \\ & \langle 2,1,6\rangle,\langle 2,3,2\rangle,\langle 2,3,5\rangle,\langle 2,5,3\rangle,\langle 2,6,5\rangle, \\ & \langle 3,1,6\rangle,\langle 3,5,2\rangle,\langle 3,5,3\rangle, \ldots,\langle 6,5,6\rangle \end{aligned}$ |
| <C, D, A*> | $\begin{aligned} & \langle 1,6,1\rangle,\langle 1,6,2\rangle,\langle 1,6,3\rangle,\langle 1,6,5\rangle,\langle 2,1,2\rangle \\ & \langle 2,6,2\rangle,\langle 2,6,3\rangle,\langle 3,2,1\rangle,\langle 3,2,3\rangle,\langle 3,2,5\rangle \\ & \langle 3,2,6\rangle,\langle 3,5,1\rangle,\langle 3,5,3\rangle, \ldots,\langle 6,5,6\rangle \end{aligned}$ |
| <D, A*, ${ }^{*}$ * | $\begin{aligned} & \langle 1,2,6\rangle,\langle 1,6,5\rangle,\langle 2,1,6\rangle,\langle 2,3,1\rangle,\langle 2,3,2\rangle \\ & \langle 2,3,5\rangle,\langle 2,5,3\rangle,\langle 2,5,6\rangle,\langle 2,6,1\rangle,\langle 2,6,5\rangle \\ & \langle 3,1,2\rangle,\langle 3,1,6\rangle,\langle 3,2,1\rangle, \ldots,\langle 6,5,6\rangle \end{aligned}$ |
| TABLE VII <br> Frequent Sequences Of 4-Itemsetsresulted From CONCATENATION PROCESS |  |
|  |  |
| <A, B, C, D> | $\begin{aligned} & \langle 1,2,1,6\rangle,\langle 1,6,3,2\rangle,\langle 1,6,3,5\rangle,\langle 1,6, \\ & 5,3\rangle,\langle 1,6,5,6>,\langle 2,1,2,6\rangle,\langle 2,1,5,3\rangle, \\ & \langle 2,1,6,1\rangle,\langle 2,3,1,6>,\langle 2,3,6,5\rangle, \ldots,\langle 6, \\ & 5,6,5\rangle \end{aligned}$ |
| <B, C, D, A*> | $\begin{aligned} & \langle 1,2,6,2\rangle,\langle 1,2,6,3\rangle,\langle 1,3,2,3\rangle,\langle 1,3 \\ & 2,5\rangle,\langle 1,5,3,5\rangle,\langle 1,5,6,5\rangle,\langle 1,6,1,6\rangle \\ & \langle 2,1,6,1\rangle,\langle 2,1,6,2\rangle,\langle 2,1,6,3\rangle, \ldots,\langle 6 \\ & 5,6,2\rangle \end{aligned}$ |
| $<\mathrm{C}, \mathrm{D}, \mathrm{A}^{*}, \mathrm{~B}^{*}>$ | $\begin{aligned} & \langle 1,6,1,6\rangle,\langle 1,6,2,1\rangle,\langle 1,6,3,2\rangle,\langle 1,6, \\ & 3,5\rangle,\langle 1,6,5,3\rangle,\langle 1,6,5,6\rangle,\langle 2,1,2,6\rangle \\ & \langle 2,6,2,1\rangle,\langle 2,6,3,1\rangle,\langle 2,6,3,6\rangle, \ldots,\langle 6, \\ & 5,6,5\rangle \end{aligned}$ |
| $\left\langle\mathrm{D}, \mathrm{A}^{*}, \mathrm{~B}^{*}, \mathrm{C}^{*}>\right.$ | $\begin{aligned} & \langle 1,2,6,5\rangle,\langle 1,6,5,3\rangle,\langle 2,1,6,3\rangle,\langle 2,1, \\ & 6,5\rangle,\langle 2,3,1,3\rangle,\langle 2,3,2,1\rangle,\langle 2,3,2,3\rangle \\ & \langle 2,3,2,5\rangle\langle 2,3,2,6\rangle,\langle 2,3,5,3>, \ldots,<6, \\ & 5,6,5\rangle \end{aligned}$ |

Next is the prune phase, where the itemsets resulting from concatenation process of 15 gendings are pruned.

The prune phase results in the function chaining of 2itemsets, 3 -itemsets, and 4 -itemsets to arrange notes.

Table VIII shows the illustration of notes sequence prunes of 2-itemsets, with note 1 used as a value in function $A$, and also as a starting point of the first note of a gatra.

The rules of note sequences are defined based on the prunes. Below is an example of rules of note sequences of gamelan music:

```
2-itemsets Rules
Start
    IF A = 1
    THEN B = 2OR 6
    ...
    IF D = 6
    THEN A* = 1 OR 2 OR 3 OR 5
```

End

## 3-itemsets Rules

 Start$$
\mathbf{I F} \mathrm{A}=1 \mathbf{A N D} \mathrm{~B}=2
$$

THEN $\mathrm{C}=1$
IF $\mathrm{C}=6 \mathrm{ANDD}=5$
THENA* $=$ 3OR6

## End

> 4-itemsets Rules Start IF A = 1 AND B = 2AND C $=1$ THEND $=6$ $\ldots$ IFD $=6$ ANDA* $^{*}=5$ AND B* $=6$ THENC* $=2$ OR5 End

TABLE VIII
Prunes Of 2-Itemsets

| A | <A, B> | <B, C> | <C, D> | <D, A*> | A |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | <1, 2> | <2, 1> | <1, 6> | <6, 1> | 1 |
|  |  |  |  | <6, 2> | 2 |
|  |  |  |  | <6, 3> | 3 |
|  |  |  |  | <6, 5> | 5 |
|  |  | <2, 3> | <3, 2> | <2, 1> | 1 |
|  |  |  |  | <2, 3> | 3 |
|  |  |  |  | <2, 5> | 5 |
|  |  |  |  | <2, 6> | 6 |
|  |  |  | <3, 5> | <5, 1> | 1 |
|  |  |  |  | <5, 2> | 2 |
|  |  |  |  | <5, 3> | 3 |
|  |  |  |  | <5, 6> | 6 |
|  |  | <2, 5> | <5, 2> | <2, 1> | 1 |
|  |  |  |  | $\langle 2,3>$ | 3 |
|  |  |  |  | <2, 5> | 5 |
|  |  |  | $\langle 5,3\rangle$ | <3, 1> | 1 |
|  |  |  |  | <3, 2> | 2 |
|  |  |  |  | <3, 5> | 5 |
|  |  |  | <5, 6> | <6, 1> | 1 |
|  |  |  |  | <6, 2> | 2 |
|  |  |  |  | <6, 3> | 3 |
|  |  |  |  | <6, 5> | 5 |

As the note sequence rules of gamelan music have been produced using AFiS algorithm, the experiment continues to melodic feature knowledge evaluation. A recommendation program to arrange note sequences is developed for the evaluation of the accuracy of the implementation of AFiS algorithm in building melodic feature knowledge of gamelan music. The knowledge and rules built by AFiS algorithm are implemented to the program. The program is expected to answer correctly some notes of test gendings randomly deleted from the sequences.

Five gendings, Liwung, Respati, Kapilaya, Tejamaya, Sengsem, are used as test gendings.

The process of data cleansing by removing pins in note sequences is implemented to all test gending, and then $30 \%$ of total notes of each test gending are changed to 0 .

Further, the note sequences test of each gending is inputted in the solving program.

The program works by scanning the notes before and after the 0 value using 2 -itemsets prunes to find the answer. The scanning process continues to 3 -itemsets prunes and 4 -itemsets prunes, and stops when there is only one note chosen as an answer. Otherwise, if there is more than 1note chosen as an answer, then the note which has the highest weight value is chosen as the answer. Table IX shows the result of the accuracy evaluation, by showing the original note sequences of each test gending, the questions, and the answers notes. The notes highlighted indicate the wrong answers by the system.

TABLE IX
Result Of The Accuracy Evaluation

| ID | TYPES | NOTE SEQUENCE |
| :---: | :---: | :---: |
| 01 | Original | $\begin{aligned} & 1,6,1,6,2,1,5,3,5,6,5,3,1,2,1,6,2,1,5,3,1 \text {, } \\ & 6,5,3,5,6,5,3,1,2,1,6 \end{aligned}$ |
|  | Question | $\begin{aligned} & 1,0,1,6,0,1,5,0,5,6,5,3,0,2,0,6,2,1,0,3,1 \text {, } \\ & 0,5,0,5,6,5,0,1,2,1,0 \end{aligned}$ |
|  | Answers | $\begin{aligned} & 1,6,1,6,3,1,5,3,5,6,5,3,1,2,1,6,2,1,5,3,1 \text {, } \\ & 6,5,3,5,6,5,3,1,2,1,6 \end{aligned}$ |
| 02 | Original | $3,2,1,6,3,2,1,6,3,2,3,5,6,3,1,6,3,5,1,6,3$, $5,1,6,5,6,5,6,3,5,3,2,6,5,3,2,6,5,3,2,3,2$, 3, 5, 6, 3, 1, 6 |
|  | Question <br> s | $3,0,1,6,0,2,0,6,3,0,3,5,6,3,1,0,3,0,1,6,0$, $5,1,0,5,6,0,6,0,5,3,2,6,5,3,0,6,5,3,0,3,2$, $3,0,6,3,1,0$ |
|  | Answers | $3,5,1,6,3,2,1,6,3,5,3,5,6,3,1,6,3,5,1,6,3$, $5,1,6,5,6,5,6,3,5,3,2,6,5,3,2,6,5,3,2,3,2$, 3, 2, 6, 3, 1, 6 |
| 03 | Original | $\begin{aligned} & 5,6,3,2,5,6,3,2,5,6,1,6,3,5,3,2,5,3,1,6,5 \\ & 3,1,6,3,5,6,5,6,5,3,2 \end{aligned}$ |
|  | Question <br> s | $\begin{aligned} & 5,0,3,2,0,6,3,0,5,0,1,0,3,5,3,0,5,3,0,6,5 \\ & 0,1,6,3,0,6,5,6,5,0,2 \end{aligned}$ |
|  | Answers | $\begin{aligned} & 5,6,3,2,3,6,3,2,5,3,1,6,3,5,3,2,5,3,5,6,5, \\ & 3,1,6,3,5,6,5,6,5,3,2 \end{aligned}$ |
| 04 | Original | $\begin{aligned} & 6,3,6,5,6,3,6,5,3,2,5,3,6,5,3,2,5,6,5,6,2, \\ & 1,5,3,2,1,2,3,2,1,6,5,6,3,6,5,6,3,6,5,3,2, \\ & 5,3,6,5,3,2,6,5,3,2,6,5,3,2,3,2,3,5,6,3,6 \\ & 5 \end{aligned}$ |
|  | Question <br> s | $\begin{aligned} & 6,3,0,5,6,3,6,0,3,0,5,3,6,0,3,0,5,6,0,6,2, \\ & 1,5,3,0,1,2,0,2,0,6,0,6,3,6,0,6,3,0,5,3,0 \\ & 5,3,6,5,0,2,6,0,3,0,6,5,3,0,3,2,0,5,6,3,0 \\ & 5 \end{aligned}$ |
|  | Answers | $\begin{aligned} & 6,3,6,5,6,3,6,5,3,2,5,3,6,5,3,2,5,6,5,6,2, \\ & 1,5,3,2,1,2,6,2,1,6,5,6,3,6,5,6,3,6,5,3,2, \\ & 5,3,6,5,3,2,6,5,3,2,6,5,3,2,3,2,3,5,6,3,6 \\ & 5 \end{aligned}$ |
| 05 | Original | $\begin{aligned} & 2,3,6,5,2,3,6,5,3,2,1,6,3,5,3,2,3,2,1,6,5 \\ & 3,5,6,1,6,5,3,2,1,6,5 \end{aligned}$ |
|  | Question <br> s | $\begin{aligned} & 2,3,6,0,2,3,6,0,3,2,1,0,3,0,3,0,3,2,1,0,5 \\ & 3,0,6,1,0,5,0,2,1,0,5 \end{aligned}$ |
|  | Answers | $\begin{aligned} & 2,3,6,5,2,3,6,5,3,2,1,6,3,5,3,2,3,2,1,6,5, \\ & 3,5,6,1,6,5,3,2,1,3,5 \end{aligned}$ |

Accuracy is measured by the $x 100 \%$ formula of (total number of right answers/total number of questions).

The results show that the accuracy of the implementation of AFiS algorithm in building melodic feature knowledge of gamelan music is up to $86.5 \%$ achieved by $(54 / 62) \times 100 \%$.

Table $X$ shows the calculation of the accuracy evaluation.

TABLE X
The Result Of Accuracy Evaluation

| The Result Of ACCURACY Evaluation |  |  |  |
| :--- | :---: | :---: | :---: |
|  | TOTAL NUMBER | TOTAL | TOTAL |
| GENDING |  | NUMBER OF | NUMBER |
|  | (TNN) | QUESTIONS | OF RIGHT |
|  | 32 | (TNN $\times$ 30\%) | ANSWERS |
| Liwung | 48 | 10 | 9 |
| Respati | 32 | 14 | 11 |
| Kapilaya | 64 | 10 | 7 |
| Tejamaya | 32 | 19 | 18 |
| Sengsem |  | 10 | 9 |
|  |  | $\mathbf{6 2}$ | $\mathbf{5 4}$ |

The second evaluation is to find whether the different answers of the program are accepted as alternative notes to the original notes. This evaluation is conducted since some sources have a difference of one or two notes in some gendings, but the difference in both sources is still accepted by gamelan artists and experts.

Two experts with an academic background and 2 experts with a practitioner background are involved to evaluate whether the different notes answered by the program can be accepted as alternative notes. Experts are asked to give a value of 0 or 1 to describe their acceptance for the different notes answered by the program, where 0 represents not accepted and 1 represents accepted. The values given by experts for each gendings are then summed up. A minimum value of 3 is used as a standard to determine that a gending can be accepted as alternative notes.

The results show that 4 out of 5 gendings, Liwung, Respati, Kapilaya, and Sengsem, are accepted by the experts.Gending Tejamaya have a value of 2 , so the different notes given by the program cannot be accepted as alternative notes. Table XI shows the evaluation of the acceptance of different note answers as alternative notes.

TABLE XI
The Result of the Evaluation of Alternative Notes

| GENDING | EXPERTS |  |  |  |  | SUM |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
|  | I ACCEPTANCE |  |  |  |  |  |
|  | II | III | IV |  |  |  |
| Liwung | 1 | 1 | 1 | 1 | 4 | Y |
| Respati | 1 | 1 | 1 | 1 | 4 | Y |
| Kapilaya | 1 | 1 | 1 | 1 | 4 | N |
| Tejamaya | 0 | 0 | 1 | 1 | 2 | N |
| Sengsem | 1 | 0 | 1 | 1 | 3 | Y |

Based on the evaluation, the different notes given by the program in 4 gendings, Liwung, Respati, and Kapilaya, and Sengsem are accepted as alternative notes, while for gending Tejamaya the different notes answered by the program are not accepted as alternative notes.

## V. Conclusion

The AFiS algorithm is a type of Apriori-based algorithm, which is suitable for small data volumes, and which was used for mining notes sequences of gamelan music. The implementation of the AFiS algorithm to build melodic feature of gamelan music results in an accuracy which is up to $86.5 \%$, while most of $13.5 \%$ that is not reached yet is accepted by the experts as alternative notes to the original notes.

Gamelan is a complex music which needs more approaches to accomplish the rules of other components of gending, such as gatra, pathet, balungan, garap.An improvement in the accuracy of the melodic feature knowledge is still needed as well as rules formulation for other components of gending. For the next research, the melodic feature knowledge resulted in this research can be used as a part of a system to develop an automatic gending composition, including the rules formulation, for a more detailed knowledge of gamelan music.

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